

A Refined Approach to Holonomic Multi-Summation

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LHCPhenonet Annual Theory Meeting 2012, Durham

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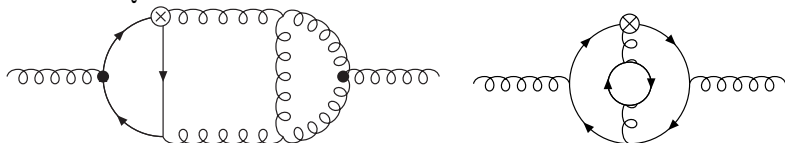
Talk Outline

- 1 A mathematical problem in particle physics
- 2 Existing tools
- 3 A refined holonomic approach
- 4 Illustrative result
- 5 Summary

High Precision QFT Calculations

QCD

- Consider perturbative QFT. Such as corrections to the high energy proton sub-structure from heavy quarks at 3-loops (and for $p^2 \gg m_{\text{HQ}}^2$). BIERENBAUM, BLÜMLEIN, KLEIN '09



- Such diagrams yield nasty integrals which are extremely hard to work with. A simpler form is needed. Our goal is to simplify the object.
- Physically speaking, the only restriction on the Feynman diagrams is that they contain at most one mass.

High Precision QFT Calculations

First Algorithmic Steps

- Momentum integrals are straightforward but the Feynman parameter integrals are not. Notice the operator insertion (the \otimes in the diagram) will give us a Mellin parameter, n . Diagrams obey homogeneous difference equations in n . BLÜMLEIN, KAUSERS, KLEIN, SCHNEIDER '09
- A first step in trying to understand the remaining integrals is to convert them into definite nested hypergeometric and harmonic multi-sums. BLÜMLEIN & KURTH '98 AND VERMASEREN '98
- There is an algorithm to do this! BLÜMLEIN, KLEIN, SCHNEIDER & STAN '12

Statement of Problem

Any given Feynman diagram (with at most one mass) can be expressed as a definite nested multi-sum,

$$F_n(\epsilon) = \sum_{i_1=\alpha_1}^{L_1(n)} \cdots \sum_{i_m=\alpha_m}^{L_m(n, i_1, \dots, i_{m-1})} \sum_{j_1=\beta_1}^{\infty} \cdots \sum_{j_N=\beta_N}^{\infty} f(\epsilon, n, i_1, \dots, i_m, j_1, \dots, j_N).$$

In this talk $\epsilon = 0$.

Aim

Given a definite nested multi-sum find, in terms of indefinite nested multi-sums, another (simpler) representation.

Statement of Problem

Illustration (< 10 minutes of CPU time)

$$\begin{aligned} & \sum_{i=2}^{n+1} \sum_{j=2}^{n+2-i} \sum_{a=0}^{\infty} \sum_{b=0}^{\infty} \frac{\binom{n+2}{i} \binom{n+2-i}{j} (-1)^{i+j} B[i, j] (j+i-1)}{(j+i+a+b)(i+j+a+b-1)(j+b)(i+b)} \\ = & \frac{2n^3 + 5n^2 + 4n - 2}{n+1} S_2 + \left[\frac{5}{2} S_2 - \frac{n^2 + 2n + 2}{n+1} \right] S_1^2 - 2n S_{2,1} \\ - & 2n(n+1) \zeta_3 - S_1^3 + \frac{1}{4} S_1^4 - \frac{1}{4} S_2^2 + 2(n+1) S_3 - \frac{1}{2} S_4 - 2S_{3,1} \\ + & [4\zeta_3 + (2n-1) S_2 + 2S_3 - 2(n+1) - 4S_{2,1}] S_1 + 2S_{2,1,1} \end{aligned}$$

where $S_{a,\dots,b} = S_{a,\dots,b}(n)$ & $B[x, y] = \Gamma[x]\Gamma[y]/\Gamma[x+y]$ for convenience.

- We have gone from a tough to simple, if long, expression.

Existing Tools

Sigma

- To solve summation problems one needs a general toolkit to apply. Of the examples available, we will use a collection of algorithms found in SCHNEIDER '04, '05... RELATED TO KARR '81 & CHYZAK '00.
- A Mathematica package, Sigma, is available that implements the algorithms.
- Sigma exploits the mathematical frame work of refined difference field theory.
- Note that Sigma computes linear recurrences with polynomial co-efficients.

Existing Tools

A Direct Approach

$$F_n = \sum_{i_1=\alpha_1}^{L_1(n)} \cdots \sum_{i_m=\alpha_m}^{L_m(n, i_1, \dots, i_m)} \sum_{j_1=\beta_1}^{\infty} \cdots \sum_{j_N=\beta_N}^{\infty} f(n, \{i\}, \{j\})$$

↓
Apply MultiSum package (tough)

$$g_0(n)F_n + \dots + g_r(n)F_{n+r} = G(n)$$

Developed and used for example by WEGSCHAIDER '08, BLÜMLEIN, KLEIN, SCHNEIDER & STAN '12

Existing Tools

Sum by Sum Approach

$$F_n = \sum_{i_1=\alpha_1}^{L_1(n)} \cdots \sum_{i_m=\alpha_m}^{L_m(n, i_1, \dots, i_m)} \sum_{j_1=\beta_1}^{\infty} \cdots \sum_{j_N=\beta_N}^{\infty} f(n, \{i\}, \{j\})$$

Zoom to the innermost sum

$$\hat{F}_n = \sum_{j_N=\beta_N}^{\infty \text{ or } L_N} f(n, \{i\}, \{j\})$$

Find a recurrence for the inner sum,
 $g_0(n)\hat{F}_n + \dots + g_r(n)\hat{F}_{n+r} = G(n)$

Solve Recurrence and
substitute in solution

Refined Holonomic Approach

$$F_n = \sum_{i_1=\alpha_1}^{L_1(n)} \cdots \sum_{i_m=\alpha_m}^{L_m(n, i_1, \dots, i_m)} \sum_{j_1=\beta_1}^{\infty} \cdots \sum_{j_N=\beta_N}^{\infty} f(n, \{i\}, \{j\})$$

Zoom to the innermost sum

$$\hat{F}_n = \sum_{j_N=\beta_N}^{\infty \text{ or } L_N} f(n, \{i\}, \{j\})$$

Find a recurrence for the inner sum,
 $g_0(n)\hat{F}_n + \dots + g_r(n)\hat{F}_{n+r} = G(n)$

Replace innermost sum
with the sequence and
its recurrences

A Worked Example

Consider the following definite nested multi-sum,

$$F_n = \sum_{i_1=0}^{n-2} \sum_{i_2=0}^{n-i_1-2} \frac{4i_1!(-1)^{i_2}(n-i_1-1)(1+i_2)!}{(1+i_1+i_2)^2(2+i_1+i_2)^2(i_1+i_2)!} \binom{n-i_1-2}{i_2}$$

To find a recurrence for F_n , first re-write the sum,

$$4 \sum_{i_1=0}^{n-2} i_1!(n-i_1-1)a_{i_1,n}$$

$$a_{i_1,n} = \sum_{i_2=0}^{n-i_1-2} \frac{(-1)^{i_2}(1+i_2)!}{(1+i_1+i_2)^2(2+i_1+i_2)^2(i_1+i_2)!} \binom{n-i_1-2}{i_2}$$

now apply a recurrence finding algorithm, e.g. Sigma.

A Worked Example

One finds that,

$$(2+i_1+n+i_1n)a_{i_1,n} - (2+i_1-n)(1+i_1+n)a_{i_1+1,n} = \frac{1}{(1+i_1^2)i_1!},$$
$$-(1+i_1-n)(1+2i_1+i_1n)a_{i,n} - (1+n)(1+i_1+i_1n)a_{i,n+1} = -\frac{1}{(1+n)i_1!}.$$

- The two recurrences plus initial conditions specify the summand — a **holonomic** sequence.
- The right-hand sides are non-zero, which is where the **'refined'** enters!

Combining them one finds, for example, a zero-order recurrence,

$$F_n = \frac{4n}{2+n} - \frac{8S(n)}{(1+n)(2+n)}, \quad S_1(n) = \sum_{i=1}^n \frac{1}{i}.$$

Illustrative Result

It is not correct to think of one method as quickest in general.
Take our example sum,

$$S = \sum_{i_1=0}^{n-2} \sum_{i_2=0}^{n-i_1-2} \frac{4i_1!(-1)^{i_2}(n-i_1-1)(1+i_2)!}{(1+i_1+i_2)^2(2+i_1+i_2)^2(i_1+i_2)!} \binom{n-i_1-2}{i_2}$$

The refined approach is 4 times faster than the current 'state-of-the-art' algorithm and some examples can be as much as 10 times faster.

Summary

- A problem in the mathematics of summation was motivated from physics.
- Various approaches and the new refined approach to multi-summation were explained.
- An illustration of the potential improvement of the algorithm was given. Although one should not think about 'best and worst' but a toolkit of algorithms.

Definition of Summand

Definition

Let $\{f(n, k)\}_{n \geq 0, k \geq 0}$ be a multivariate (here bivariate) sequence over a field \mathbb{F} . f is hypergeometric in n and k if $\exists d \in \mathbb{Z}^+$ and $r_1(x, y), r_2(x, y) \in \mathbb{F}(x, y)$ such that $f(n+1, k)/f(n, k) = r_1(n, k)$ and $f(n, k+1)/f(n, k) = r_2(n, k) \forall n, k \geq d$.

Definition

Restrict to sequences where:

$\exists d \in \mathbb{Z}^+$, $a_i, b_i, m_i \in \mathbb{Z}$ and $c_i \in \mathbb{F}$ for $i \in [1, r]$ where $a_i n + b_i k + c_i \notin \mathbb{Z}^- \forall n, k \geq d$ and $\exists p(x, y), q(x, y) \in \mathbb{F}[x, y]$ where $q \neq 0$ factors linearly over \mathbb{F} :

$$f(n, k) = p(n, k)/q(n, k) \prod_{i=1}^r \Gamma(a_i n + b_i k + c_i)^{m_i} \forall n, k \geq d.$$

n.b. \mathbb{Z}^\pm both include zero.

Statement of Problem

Summary

Allow the summand to be as defined, times harmonic numbers and linear combinations of such terms.

$$F_n(\epsilon) = \sum_{i_1=\alpha_1}^{L_1(n)} \cdots \sum_{i_m=\alpha_m}^{L_m(n, i_1, \dots, i_m)} \sum_{j_1=\beta_1}^{\infty} \cdots \sum_{j_N=\beta_N}^{\infty} f(\epsilon, n, i_1, \dots, i_m, j_1, \dots, j_N)$$

where $n \in \mathbb{N}$ and $n \geq n_0 \in \mathbb{N}$, all the upper bounds, L_k , are integer-linear in n , $\{i\}$ and $\{j\}$. All the lower bounds are specified constants, $\{\alpha\}, \{\beta\} \in \mathbb{N}$. The summand, F , is hypergeometric (defined shortly) with respect to n , $\{i\}$ and $\{j\}$. In this talk $\epsilon = 0$

Aim

Find an expression for a definite nested multi-sum, using the outlined summand, in terms of indefinite nested multi-sums.