Reduction of one and two-loop amplitudes at the Integrand level

◆□▶ ◆□▶ ▲□▶ ▲□▶ ▲□ ◆ ○ ◆ ○ ◆

Who?Ioannis MalamosFrom?Institute de Fisica Corpuscular,
ValenciaWhen?Durham, 22 March 2012

Introduction

- Phenomenology at new colliders requires NLO (and NNLO) calculations
- One of the parts of these calculations is the virtual part (loop diagrams)
- Reduction methods are used due to the complexity and the large number of Feynman integrals
- One-loop case completely solved, towards a two-loop reduction method

Historical Background

- Already attempts in the 60's (D.B.Melrose (1965),G.Källèn and J.Toll (1965))
- Passarino-Veltman reduction (general applicability, major achievements but not designed at amplitude level)
- Unitarity based methods-Bern,Dixon,Dunbar,Kosower (major advantage: designed to work at amplitude level, limited applications)
- Quadruple and triple cuts- Britto, Cachazo, Feng (major simplifications)

 Reduction at the integrand level (Ossola,Papadopoulos,Pittau)

Definitions

A general scalar Feynman integral of order n is given

$$\int d^4q \frac{1}{D_1 D_2 \dots D_n} \tag{1}$$

with $D_i = (q + p_i)^2 - m_i^2 = q^2 + 2p_i \cdot q + \mu_i$ the inverse propagators

p_i's are the momenta entering the propagators, related to the external momenta

we don't assume momentum conservation, we rather deal with integrand-Graphs or i-Graphs given by

$$\frac{1}{D_1 D_2 \dots D_n}$$

instead of Feynman Graphs We deal with scalar integrals without loss of generality

Reduction at one loop

Write the number 'one' in terms of denominators

$$1 = \sum_{i}^{n} T_{i}(q) D_{i}$$
(2)

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● のへで

If we find such T_i 's then reduction is achieved

$$\frac{1}{D_1 \cdots D_n} = \frac{T_1(q)}{D_2 D_3 \cdots D_n} + \dots + \frac{T_n(q)}{D_1 D_2 \cdots D_{n-1}} ,$$
(3)

Equation polynomial in q

Reduction with trivial coefficients

We start by assuming that the T_i 's are constants in the loop momentum

$$T_j(q)=x_j$$

For any value of q

$$q^{2}\sum_{j=1}^{n}x_{j}+2q_{\mu}\sum_{j=1}^{n}x_{j}p_{j}^{\mu}+\sum_{j+1}^{n}x_{j}\mu_{j}=1$$
 (4)

For the d + 2 equations we need d + 2 coefficients In 4 dimensions a hexagon is decomposed to pentagons with trivial coefficients Reduction with coefficients linear in q

We assume now

$$T_j(q) = x_j + \sum_{k=1}^4 x_{j,k}(q \cdot t_k)$$
 (5)

• t_k are d linearly independent arbitrary vectors, forming a base in d dimensions

- We start with $(d + 1) \times n$ coefficients, not all being independent.
- The tensor structures we have to construct are denoted by

1 ,
$$q^{\mu}$$
 , $q^{\mu}q^{\nu}$, $q^{2}q^{\mu}$. (6)

Tensor structures are given by $\frac{d^2+5d}{2}$ + 1 (19 in d = 4)

Number of independent coefficients numerically (with rounding errors)

n	<i>d</i> = 6	<i>d</i> = 5	<i>d</i> = 4	<i>d</i> = 3	<i>d</i> = 2	<i>d</i> = 1
2	14-0	12-0	10-0	8-0	6-0	4-0
3	21-1	18-1	15-1	12-1	9-1	6-2
4	28 -3	24-3	20-3	16-3	12-4	8-4
5	35-6	30-6	25-6	20-7	15-7	10-6
6	42-10	36-10	30-11	24-11	18-10	12-8
7	49-15	42-16	35-16	28-15	21-13	14-10
8	56-22	48-22	40-21	32-19	24-16	16-12

The rank of *M* for various *d* and *n*, given as the difference $\xi - q$.

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● のへで

Decomposability with linear terms

- Every integral of order d + 1 is decomposed to an integral of order d with coefficients linear in the loop momentum
- In the most interesting case of d = 4 a pentagon is decomposed to boxes
- Situation does not improve with coefficients of higher order

Uniqueness of the decomposition

 A generic two loop diagram has three kind of propagators

$$D(l_{1} + p_{1}), D(l_{1} + p_{2}), ..., D(l_{1} + p_{n}),$$

$$D(l_{2} + q_{1}), D(l_{2} + q_{2}), ..., D(l_{2} + q_{k}),$$

$$D(l_{1} + l_{2} + r_{1}), D(l_{1} + l_{2} + r_{2}), ..., D(l_{1} + l_{2} + r_{s}),$$

(7)

- We denote this iGraph by (n,s,k) and we define it to be of order n + k + s.
- Due to these mixed propagators the problem is not just a double copy of a one-loop case

Reduction at two loops

- At one-loop the base of Master Integrals is known a priori.At two loop not
- Cannot be just scalar integrals (i.e. $l_1.q_1$ is not reducible, notice the difference with one loop)
- Reduction seems to depend on what kind of iGraphs one wants to reduce to

- Our motivation is to find out what is the highest number of denominators every iGraph can be decomposed to (i.e. in one-loop it is d)
- For unitarity reasons we expect this to be 2d

Reduction at two loops

'Counting to one' again

$$\sum_{j=1}^{n_1} x_j D(l_1 + p_j) + \sum_{j=n_1+1}^{n_1+n_2} x_j D(l_2 + p_j) + \sum_{j=n_1+n_2+1}^{n} x_j D(l_1 + l_2 + p_j) = 1$$
(8)

< □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □

Notice that $n_{1,2,3} \leq d$, otherwise the problem can be solved as a one loop problem

Reduction at two loops

- Imagine the iGraph (5,1,1) in 4 dimensions. We can put all coefficients of l_2 to zero and solve the pentagon to boxes problem for l_1 only (as we saw before with linear terms)
- Restricted number of iGraphs to analyse-highest being the (d,d,d)
- An iGraph of order 2d + 4 is decomposed to an iGraph of order 2d + 3 with trivial coefficients x_i

(日) (日) (日) (日) (日) (日) (日)

Reduction at two loops with linear terms

We try now the linear terms of the following type

$$x_{j} = \sum_{i} (a_{j} + b_{ij}(l_{1} \cdot t_{i}) + c_{ij}(l_{2} \cdot t_{i}))$$
(9)

(日) (日) (日) (日) (日) (日) (日)

We again find the number of independent coefficients (with rounding errors) and compare it with the number of tensor structures. We give the table with our findings, using again a horizontal line for the reducible cases. T(d) is the number of tensor structures Reduction at two loops with linear terms

$$L = 2$$
, linear
 $T(d) = (4d^2 + 18d + 2)/2$

n	<i>d</i> = 6	<i>d</i> = 5	<i>d</i> = 4	<i>d</i> = 3	<i>d</i> = 2
3	39-0	33-0	27-0	21-0	15-0
4	52-0	44-0	36-0	28-0	20-0
5	65-1	55-1	45-1	35-1	25-1
6	78-3	66-3	54-3	42-3	30-3
7	91-6	77-6	63-6	49-6	35-8
8	104-10	88-10	72-10	56-10	40-10
9	111-15	99-15	81-15	63-17	45-18
10	130-21	110-21	90-21	70-24	50-23
11	143-28	121-28	99-30	77-31	55-28
12	156-36	132-36	108-39	84-36	60-33
13	169-45	143-47	117-48	91-45	65-38
14	182-55	154-58	126-57	98-52	70-43
T (d)	127	96	69	46	27

Reduction at two loops with linear terms

In d dimensions, every 2d + 2 iGraph is decomposed to a 2d + 1 with linear terms.

 In order to go one step further we have to consider 2d + 1 iGraphs with coefficients of quadratic,cubic,quartic dependence in the loop momenta

The set of iGraphs of interest is even more restricted. In d = 2 we have to consider the (2,1,2) only, in d = 3 the (3,1,3),(3,2,2) and in d = 4 the (4,1,4),(4,3,2) and (3,3,3)

(日) (日) (日) (日) (日) (日) (日)

Reduction at two loops with higher order terms

- We repeat the proceedure using quadratic, and cubic terms terms. The reducibility of diagrams with higher terms depends on the number of dimensions
- We find that the (2,1,2) is reducible in 2 dimensions with quadratic terms, while the similar 2d + 1 iGraphs in 3 and 4 dimensions are not.

- We find solutions for the iGraphs of order 7 in 3 dimensions with cubic terms
- We find solutions for the iGraphs of order 9 in 4 dimensions with cubic terms

Reduction at two loops with higher order terms

- For all the cases above we solve numerically the 1 = 1 equation.
- We construct a base for two loop iGraphs with at most 2d denominators.We proved it in 2,3 and 4 dimensions.
- We are interested in a base for the moment (not necessarily the minimal base)
- Once again, when considering iGraphs with 2d denominators, they must have a maximal cut. For example, a (6,1,1) in 4 dimensions does not belong in this category as we saw and as expected from Unitarity.

Not the end of the story yet!

 Consider the following Feynman integral in 2 dimensions

$$\int d^2 l_1 d^2 l_2 \frac{1}{l_1^2 (l_1 + p)^2 (l_1 + l_2)^2 l_2^2 (l_2 - p)^2}$$
(10)

It should be decomposable. We write:

$$1 = a_1 l_1^2 + a_2 (l_1 + p)^2 + \dots + a_5 (l_2 - p)^2 \qquad (11)$$

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

There should be coefficients that this holds for any l_1 and l_2

Not the end of the story yet!

Find l_1 and l_2 such that

$$l_1^2 = (l_1 + p)^2 = l_2^2 = (l_2 - p)^2 = 0$$
 (12)

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ ▲ □ ● ● ● ●

However, for such I_1 and I_2 also

 $(I_1 + I_2)^2 = 0$

It means that there are l_1 and l_2 for which 1 = 0! This diagram is NOT decomposable!

Solution

- The diagram above has a problemtaic maximal cut. However, there is a category of terms that one can add to the 1 = 1 equation that vanish upon integration and can make the equation hold for every l_1 and l_2 .
- Any total derivative in dimensional regularisation vanishes (I.B.P) (Chetyrkin and Tkachov)

An example of an I.B.P. identity is the following

$$\int \frac{\partial}{\partial l_1^{\mu}} \left(\frac{(l_1 + p_1)^{\mu}}{D(l_1 + p_1)D(l_2 + p_2)D(l_1 + l_2 + p_3)} \right) = 0$$
(13)

◆□▶ ◆□▶ ▲□▶ ▲□▶ ▲□ ◆ ○○

Solution

 One can now add at the Integrand level all these total derivatives (together with the previous terms we had) and ask for coefficients

The problem is now solved!

For the particular 2-dimensional problem, after integrating we get the following solution (which agrees with the literature)

$$\int \frac{1}{l_1^2(l_1+p)^2(l_1+l_2)^2l_2^2(l_2-p)^2} = \frac{-1}{p^2} \int \frac{1}{l_1^2(l_1+p)^2l_2^2(l_2-p)^2} + \frac{4}{(p^2)^2} \int \frac{1}{l_1^2(l_1+l_2)^2(l_2-p)^2}$$

(***) ののの 手 〈言〉〈手〉〈馬〉〈与〉

I.B.P. at the Integrand level

We tested the I.B.P.'S equations for other known examples and they work

As a by-product, we can use the I.B.P.'s at the integrand level at one loop as well and generalise the OPP method with denominators of higher powers

(日) (日) (日) (日) (日) (日) (日)

Conclusions

- We presented Reductions at the Integrand level for one and two loop amplitudes
- Although the one loop case is already known we could rediscover things with a slightly different method. Important because the coefficients of the reduction depend on the base one asks to reduce to
- At two loops we showed why we expect unitarity to work, and how by simple counting one can decompose two-loop iGraphs
- For special cases one needs to use Integration By Parts identities. We use them at the integrand level for the first time

Conclusions

In principle both counting and I.B.P's should be used
 Our method is very general (does not depend on masses or planar or non-planar Graphs) and is very simple as well.

◆□▶ ◆□▶ ▲□▶ ▲□▶ ▲□ ◆ ○ ◆ ○ ◆

BACKUP SLIDES

◆□▶ ◆□▶ ◆ □▶ ◆ □▶ ● □ ● ● ● ●

Finding the independent coefficients

- Choose random values for external momenta and masses
- Choose as many random values for q as the number of coefficients and substitute in equation 2
- For these ξ values we get a set of ξ linear equations for the ξ unknowns x:

$$\sum_{j=1}^{\xi} M^{i}_{\ j} x^{j} = 1 \quad , \quad j = 1, \dots, \xi \quad . \tag{15}$$

In case the number of independent tensor structures we can form are less than ξ the determinant of M vanishes. In a numerical computation there will be of course rounding errors

Finding the independent coefficients

- For q zero eigenvalues this Matrix will have a determinant of the order 10^{-pq} where p is the number of significant digits
- Running for different precisions we get the value q and the rank of the Matrix from ξq
- If the rank of the Matrix is equal to the number of tensor structures the integral is decomposable
- For different dimensions and number of propagators we give the results of the number of independent coefficients
- We denote the limit of decomposability with horizontal lines

Decomposability with spurious terms

 In the conventional approach, a pentagon is decomposed is decomposed to boxes with spurious terms

$$1 = \sum_{i=1}^{5} (a_i + \tilde{a}_i S_i(q)) D_i$$
 (16)

- The S_i 's vanish after integration, by construction. Their form (q-dependence) is known
- Triangles in our case vanish giving the same decomposition
- The resulting base for one loop integrals is scalar integrals up to 4 denominators