# Multiloop Integrand Reduction via Multivariate Polynomial Division

#### Tiziano Peraro

Max-Planck-Institut für Physik Föhringer Ring 6, D-80805 München, Germany

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#### T. Peraro (MPI - München) Multiloop Integrand Reduction via Multivariate Polynomial Division

RADCOR2013

#### Based on:

- P. Mastrolia, E. Mirabella and **T.P.**, *Integrand reduction of one-loop scattering amplitudes through Laurent series expansion*, JHEP **1206**, 095 (2012) [arXiv:1203.0291 [hep-ph]].



P. Mastrolia, E. Mirabella, G. Ossola and **T.P.**, *Scattering Amplitudes from Multivariate Polynomial Division*, Phys. Lett. B **718**, 173 (2012) [arXiv:1205.7087 [hep-ph]].



P. Mastrolia, E. Mirabella, G. Ossola and **T.P.**, *Integrand-Reduction for Two-Loop Scattering Amplitudes through Multivariate Polynomial Division*, Phys. Rev. D 87, 085026 (2013) [arXiv:1209.4319 [hep-ph]].



H. van Deurzen, N. Greiner, G. Luisoni, P. Mastrolia, E. Mirabella, G. Ossola, **T.P.**, J. F. von Soden-Fraunhofen and F. Tramontano Phys. Lett. B **721**, 74 (2013) [arXiv:1301.0493 [hep-ph]].





G. Cullen, H. van Deurzen, N. Greiner, G. Luisoni, P. Mastrolia, E. Mirabella, G. Ossola, **T.P.** and F. Tramontano *NLO QCD* corrections to Higgs boson production plus three jets in gluon fusion, arXiv:1307.4737 [hep-ph].



P. Mastrolia, E. Mirabella, G. Ossola and T.P., Multiloop Integrand Reduction for Dimensionally Regulated Amplitudes, arXiv:1307.5832 [hep-ph].



H. van Deurzen, G. Luisoni, P. Mastrolia, E. Mirabella, G. Ossola and **T.P.**, *NLO QCD corrections to Higgs boson production in association with a top quark pair and a jet*, arXiv:1307.8437 [hep-ph].

#### Outline



- Introduction and motivation
- 2 The integrand reduction of scattering amplitudes
- Integrand reduction via polynomial division
  - Application at one-loop
- 5 Higher loops
- 6 Conclusions

# Introduction and motivation

#### Motivation

- Understanding the basic analytic and algebraic structure of integrands and integrals of scattering amplitudes
- Exploration of methods for obtaining theoretical predictions in perturbative Quantum Field Theory at higher orders
- Automation of the computation of loop integrals

We developed a coherent framework for the integrand decomposition of Feynman integrals

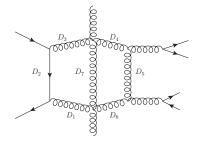
- based on simple concepts of algebraic geometry
- applicable at all loops

#### Integrand reduction

- Generic *l*-loop integral:
  - is a rational function in the components of the loop momenta  $q_i$
  - polynomial numerator  $\mathcal{N}_{i_1...i_n}$

$$\mathcal{M}_n = \int d^d q_1 \dots d^d q_\ell ~~ \mathcal{I}_{i_1 \dots i_n}, \qquad \mathcal{I}_{i_1 \dots i_n} \equiv rac{\mathcal{N}_{i_1 \dots i_n}}{D_{i_1} \dots D_{i_n}}$$

- quadratic polynomial denominators D<sub>i</sub>
  - they correspond to Feynman loop propagators



$$D_i = \left(\sum_j (-)^{s_{ij}} q_j + p_i\right)^2 - m_i^2$$

 $\dots D_i$ 

2

# Integrand reduction

#### The idea

Manipulate the integrand and reduce it to a linear combination of "simpler" integrands.

• The integrand-reduction algorithm leads to

$$\mathcal{I}_{i_1\cdots i_n} \equiv \frac{\mathcal{N}_{i_1\cdots i_n}}{D_{i_1}\cdots D_{i_n}} = \frac{\Delta_{i_1\cdots i_n}}{D_{i_1}\cdots D_{i_n}} + \ldots + \sum_{k=1}^n \frac{\Delta_{i_k}}{D_{i_k}} + \Delta_{\emptyset}$$

• The residues  $\Delta_{i_1...i_k}$  are irreducible polynomials in  $q_i$ 

- can't be written as a combination of denominators  $D_{i_1} \dots D_{i_k}$
- universal topology-dependent parametric form
- the coefficients of the parametrization are process-dependent

# From integrands to integrals

• By integrating the integrand decomposition

$$\mathcal{M}_n = \int d^d q_1 \dots d^d q_\ell \left( rac{\Delta_{i_1 \dots i_n}}{D_{i_1} \dots D_{i_n}} + \dots + \sum_{k=1}^n rac{\Delta_{i_k}}{D_{i_k}} + \Delta_{\emptyset} 
ight)$$

- some terms vanish and do not contribute to the amplitude ⇒ spurious terms
- non-vanishing terms give Master Integrals (MIs)
- The amplitude is a linear combination of MIs
- The coefficients of this linear combination can be identified with some of the coefficients which parametrize the polynomial residues

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 $\Rightarrow$  reduction to MIs  $\equiv$  polynomial fit of the residues

#### The one-loop decomposition

At one loop the result is well known:

• the integrand decomposition [Ossola, Papadopoulos, Pittau (2007); Ellis, Giele, Kunszt, Melnikov (2008)]

$$\begin{aligned} \mathcal{I}_{i_1 \cdots i_n} &= \frac{\mathcal{N}_{i_1 \cdots i_n}}{D_{i_1} \cdots D_{i_n}} = \sum_{j_1 \dots j_5} \frac{\Delta_{j_1 j_2 j_3 j_4 j_5}}{D_{j_1} D_{j_2} D_{j_3} D_{j_4} D_{j_5}} + \sum_{j_1 j_2 j_3 j_4} \frac{\Delta_{j_1 j_2 j_3 j_4}}{D_{j_1} D_{j_2} D_{j_3} D_{j_4}} \\ &+ \sum_{j_1 j_2 j_3} \frac{\Delta_{j_1 j_2 j_3}}{D_{j_1} D_{j_2} D_{j_3}} + \sum_{j_1 j_2} \frac{\Delta_{j_1 j_2}}{D_{j_1} D_{j_2}} + \sum_{j_1} \frac{\Delta_{j_1}}{D_{j_1}} \end{aligned}$$

the integral decomposition

$$= c_{4,0} + c_{3,0} + c_{2,0} + c_{1,0} + c_{1,0} + c_{4,4} + c_{3,7} + c_{2,9} + c_$$

5

### Integrand reduction and polynomials

• At *l*-loops we want to achieve the integrand decomposition:

$$\mathcal{I}_{i_1\dots i_n}(q_1,\dots,q_\ell) \equiv \frac{\mathcal{N}_{i_1\dots i_n}}{D_{i_1}\dots D_{i_n}} = \underbrace{\frac{\Delta_{i_1\dots i_n}}{D_{i_1}\dots D_{i_n}}}_{\text{they must be irreducible}} + \dots + \sum_{k=1}^n \frac{\Delta_{i_k}}{D_{i_k}} + \Delta_{\emptyset}$$

We trade (q<sub>1</sub>,...,q<sub>ℓ</sub>) with their coordinates z ≡ (z<sub>1</sub>,..., z<sub>m</sub>)
 ⇒ numerator and denominators ≡ polynomials in z

$$\mathcal{I}_{i_1\ldots i_n}(\mathbf{z})\equiv rac{\mathcal{N}_{i_1\ldots i_n}(\mathbf{z})}{D_{i_1}(\mathbf{z})\ldots D_{i_n}(\mathbf{z})}$$

 $\Rightarrow$  Integrand reduction  $\equiv$  problem of multivariate polynomial division

The problem of the determination of the residues of a generic diagramhas been solved.[Y. Zhang (2012), P. Mastrolia, E. Mirabella, G. Ossola, T.P. (2012)]

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#### Residues via polynomial division

- Y. Zhang (2012), P. Mastrolia, E. Mirabella, G. Ossola, T.P. (2012)
  - Define the Ideal of polynomials

$$\mathcal{J}_{i_1\cdots i_n} \equiv \langle D_{i_1},\ldots,D_{i_n} \rangle = \left\{ p(\mathbf{z}) \, : \, p(\mathbf{z}) = \sum_j h_j(\mathbf{z}) D_j(\mathbf{z}), \, h_j \in P[\mathbf{z}] \right\}$$

• Take a Gröbner basis 
$$G_{\mathcal{J}_{i_1\cdots i_n}}$$
 of  $\mathcal{J}_{i_1\cdots i_n}$ 

$$G_{\mathcal{J}_{i_1\cdots i_n}} = \{g_1, \dots, g_s\}$$
 such that  $\mathcal{J}_{i_1\cdots i_n} = \langle g_1, \dots, g_s \rangle$ 

• Perform the multivariate polynomial division  $\mathcal{N}_{i_1...i_n}/G_{\mathcal{J}_{i_1...i_n}}$ 

$$\mathcal{N}_{i_1\cdots i_n}(z) = \underbrace{\sum_{k=1}^n \mathcal{N}_{i_1\cdots i_{k-1}i_{k+1}\cdots i_n}(z) D_{i_k}(z)}_{\text{quotient} \in \mathcal{J}_{i_1\cdots i_n}} + \underbrace{\Delta_{i_1\cdots i_n}(z)}_{\text{remainder}}$$

• The remainder  $\Delta_{i_1 \cdots i_n}$  is irreducible  $\Rightarrow$  can be identified with the residue

# Recursive Relation for the integrand decomposition

P. Mastrolia, E. Mirabella, G. Ossola, T.P. (2012)

#### The recursive formula

$$\mathcal{N}_{i_1\cdots i_n} = \sum_{k=1}^n \mathcal{N}_{i_1\cdots i_{k-1}i_{k+1}\cdots i_n} D_{i_k} + \Delta_{i_1\cdots i_n}$$
 $\mathcal{I}_{i_1\cdots i_n} \equiv rac{\mathcal{N}_{i_1\cdots i_n}}{D_{i_1}\cdots D_{i_n}} = \sum_k \mathcal{I}_{i_1\cdots i_{k-1}i_{k+1}\cdots i_n} + rac{\Delta_{i_1\cdots i_n}}{D_{i_1}\cdots D_{i_n}}$ 

- Fit-on-the-cut approach
  - from a generic  ${\cal N},$  get the parametric form of the residues  $\Delta$
  - determine the coefficients sampling on the cuts (impose  $D_i = 0$ )
- Divide-and-Conquer approach
  - $\bullet\,$  generate the  ${\cal N}$  of the process
  - compute the residues by iterating the polynomial division algorithm

[Ossola, Papadopoulos, Pittau (2007)]

The decomposition of the numerator

$$\mathcal{N}_{i_1\cdots i_n} = \Delta_{i_1\cdots i_n} + \sum_k \Delta_{i_1\cdots i_{k-1}i_{k+1}\cdots i_n} D_{i_k} + \dots$$

- Fit the coefficients of the residues sampling on the multiple cuts
- First step: n-ple cut

• impose 
$$D_{i_1} = \ldots = D_{i_n} = 0$$

$$\Delta_{i_1\cdots i_n}=\mathcal{N}_{i_1\cdots i_n}$$

- Further steps: k-ple cut
  - impose  $D_{i_1} = \ldots = D_{i_k} = 0$  for any subset  $\{i_1 \ldots i_k\}$

$$\Delta_{i_1\cdots i_k} = \frac{\mathcal{N}_{i_1\cdots i_n} - \text{higher-point contibutions}}{\prod_{h\neq i_1,\dots,i_k} D_h}$$

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# Fit-on-the-cut approach: The reducibility criterion

What happens if a cut has no solution?

#### The reducibility criterion

- If a cut D<sub>i1</sub> = ... = D<sub>ik</sub> = 0 has no solutions, the associated residue vanishes. In other words, any numerator is completely reducible.
- This generally happens with overdetermined systems i.e. when the number of cut denominators is higher than the one of loop coordinates.

• When 
$$D_{i_1} = \ldots = D_{i_k} = 0$$
 has no solution:

 $\Delta_{i_1...i_k} = 0 \implies \text{no need to perform the fit}$  $\mathcal{N}_{i_1...i_n} = \sum_{k=1}^n \mathcal{N}_{i_1...i_{k-1}i_{k+1}...i_n} D_{i_k}$  $\mathcal{I}_{i_1...i_n} = \sum_k \mathcal{I}_{i_1...i_{k-1}i_{k+1}...i_n}$ 

### Fit-on-the-cut approach: The maximum-cut theorem

#### The maximum-cut theorem

We define maximum-cut, a cut where

 $#(cut-denominators) \equiv #(components-of-loop-momenta)$ 

In non-special kinematic configurations it has a finite number of solutions

#(coefficients-of-the-residue) = #(solutions-of-the-cut)

• The fit-on-the-cut approach therefore gives a number of equations which is equal to the number of unknown coefficients.

# Fit-on-the-cut approach: The maximum-cut theorem

#### Examples:

diagram	Δ	$n_s$	diagram	Δ	$n_s$
$\langle \downarrow \rangle$	$c_0$	1	Ц	$c_0 + c_1 z$	2
$\langle \square$	$\sum_{i=0}^{3} c_i z^i$	4	$\langle \times$	$\sum_{i=0}^{3} c_i z^i$	4
E	$\sum_{i=0}^{7} c_i z^i$	8		$\succ \sum_{i=0}^{7} c_i z^i$	8

Pros:

- each multiple cut projects out the corresponding residue
  - $\Rightarrow$  the systems of equations for the coefficients are much smaller
- can be implemented either analytically or numerically
- very successful application at one-loop

Cons:

- at higher-loops the solutions of the cuts can be difficult to find
- it cannot be applied in the presence of higher powers of denominators
  - a cut denominator might be equal to an uncut denominator

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  - it that case

$$\frac{\mathcal{N}_{i_1\cdots i_n} - \text{higher-point contibutions}}{\prod_{h \neq i_1, \dots, i_k} D_h} = \frac{0}{0}$$

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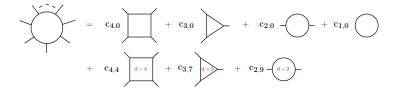
Cons:



#### One-loop decomposition from polynomial division

P. Mastrolia, E. Mirabella, G. Ossola, T.P. (2012)

- Start from the most general one-loop amplitude in  $d = 4 2\epsilon$
- Apply the recursive formula for the integrand decomposition
  - ⇒ it reproduces the OPP result [Ossola, Papadopoulos, Pittau (2007); Ellis, Giele, Kunszt, Melnikov (2008)]
- Drop the spurious terms
- ⇒ Get the most general integral decomposition (well knwon result)



### One-loop decomposition from polynomial division

At one loop in  $4 - 2\epsilon$  dimensions:

- 5 coordinates  $\mathbf{z} = (z_1, z_2, z_3, z_4, z_5)$ 
  - 4 components  $(z_1, z_2, z_3, z_4)$  of q w.r.t. a 4-dimensional basis
  - $z_5 = \mu^2$  encodes the  $(-2\epsilon)$ -dependence on the loop momentum

we start with

$$\mathcal{I}_n \equiv \mathcal{I}_{1...n} = \frac{\mathcal{N}_{1...n}(\mathbf{z})}{D_1(\mathbf{z}) \dots D_n(\mathbf{z})} \qquad \text{most general 1-loop numerator}$$
generic 1-loop denominators

• if m > 5 any integrand  $\mathcal{I}_{i_1...i_m}$  is reducible (reducibility criterion)

$$\mathcal{I}_{i_1\cdots i_m} = \sum_k \mathcal{I}_{i_1\cdots i_{k-1}i_{k+1}\cdots i_m}, \quad \Rightarrow \quad \Delta_{i_1\cdots i_m} = 0 \quad \text{for } m > 5$$

 for *m* ≤ 5 the polynomial-division algorithm gives the already-known parametric form of the residues Δ<sub>ijk...</sub> Choice of 4-dimensional basis for an *m*-point residue

$$e_1^2 = e_2^2 = 0$$
,  $e_1 \cdot e_2 = 1$ ,  $e_3^2 = e_4^2 = \delta_{m4}$ ,  $e_3 \cdot e_4 = -(1 - \delta_{m4})$ 

• Coordinates:  $\mathbf{z} = (z_1, z_2, z_3, z_4, z_5) \equiv (x_1, x_2, x_3, x_4, \mu^2)$ 

$$q_{4-\text{dim}}^{\mu} = -p_{i_1}^{\mu} + x_1 \ e_1^{\mu} + x_2 \ e_2^{\mu} + x_3 \ e_3^{\mu} + x_4 \ e_4^{\mu}, \qquad q^2 = q_{4-\text{dim}}^2 - \mu^2$$

Generic numerator

$$\mathcal{N}_{i_1\cdots i_m} = \sum_{j_1,\cdots,j_5} \alpha_{\vec{j}} \, z_1^{j_1} \, z_2^{j_2} \, z_3^{j_3} \, z_4^{j_4} \, z_5^{j_5}, \qquad (j_1 \dots j_5) \quad \text{such that} \quad \operatorname{rank}(\mathcal{N}_{i_1 \dots i_m}) \le m$$

Residues

$$\begin{aligned} \Delta_{i_1 i_2 i_3 i_4 i_5} &= c_0 \\ \Delta_{i_1 i_2 i_3 i_4} &= c_0 + c_1 x_4 + \mu^2 (c_2 + c_3 x_4 + \mu^2 c_4) \\ \Delta_{i_1 i_2 i_3} &= c_0 + c_1 x_3 + c_2 x_3^2 + c_3 x_3^3 + c_4 x_4 + c_5 x_4^2 + c_6 x_4^3 + \mu^2 (c_7 + c_8 x_3 + c_9 x_4) \\ \Delta_{i_1 i_2} &= c_0 + c_1 x_2 + c_2 x_3 + c_3 x_4 + c_4 x_2^2 + c_5 x_3^2 + c_6 x_4^2 + c_7 x_2 x_3 + c_9 x_2 x_4 + c_9 \mu^2 \\ \Delta_{i_1} &= c_0 + c_1 x_1 + c_2 x_2 + c_3 x_3 + c_4 x_4 \end{aligned}$$

#### It can be easily extended to higher-rank numerators

#### Fit-on-the-cut at 1-loop

# Integrand decomposition: $=\Sigma$ Fit-on-the cut fit *m*-point residues on *m*-ple cuts $=\Sigma + \Sigma + \Sigma + \Sigma + \Sigma + \Sigma + \Sigma + (1 + 2 + 2) + (1 + 2) + ($

#### Improved 1-loop Reduction with NINJA

P. Mastrolia, E. Mirabella, T.P. (2012)

The integrand reduction via Laurent expansion:

- fits residues by taking their asymptotic expansions on the cuts
- yields diagonal systems of equations for the coefficients
- requires the computation of fewer coefficients
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- ★ Implemented in the semi-numerical C++ library NINJA
  - Laurent expansions via a simplified polynomial-division algorithm
  - interfaced with the package GOSAM
  - is a faster and more stable integrand-reduction algorithm

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  - is a faster and more stable integrand-reduction algorithm
- $\Rightarrow$  see P. Mastrolia's talk for more details

#### Extension to higher loops

- The integrand-level approach to scattering amplitudes at one-loop
  - can be used to compute any amplitude in any QFT
  - has been implemented in several codes, some of which public [SAMURAI, CUTTOOLS, NGLUONS]
  - has produced (and is still producing) results for LHC [GOSAM (see P. Mastrolia's talk),

FORMCALC, BLACKHAT, MADLOOP, NJETS, OPENLOOP ...]

- At two or higher loops
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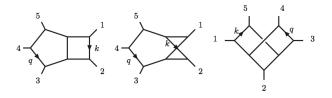
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The integrand-level approach might be a tool for understanding the structure of multi-loop scattering amplitudes and a method for their evaluation.

• ... we are moving the first steps in this direction

# $\mathcal{N}=4$ SYM and $\mathcal{N}=8$ SUGRA amplitudes

P. Mastrolia, G. Ossola (2011); P. Mastrolia, E. Mirabella, G. Ossola, T.P. (2012)



- Examples in  $\mathcal{N} = 4$  SYM and  $\mathcal{N} = 8$  SUGRA amplitudes (d = 4)
  - generation of the integrand
    - graph based [Carrasco, Johansson (2011)]
    - unitarity based [U. Schubert (Diplomarbeit)]
  - fit-on-the-cut approach for the reduction
- Results:
- $\mathcal{N}=4~$  linear combination of 8 and 7-denominators MIs
- $\mathcal{N}=8$  linear combination of 8, 7 and 6-denominators MIs

### Divide-and-Conquer approach

P. Mastrolia, E. Mirabella, G. Ossola, T.P. (2013)

The divide-and-conquer approach to the integrand reduction

- does not require the knowledge of the solutions of the cut
- can always be used to perform the reduction in a finite number of purely algebraic operations
- has been automated in a PYTHON package which uses MACAULAY2 and FORM for algebraic operations







 also works in special cases where the fit-on-the-cut approach is not applicable (e.g. in presence of double denominators)

### Divide-and-Conquer approach: a simple example

$$\begin{array}{cccc}
\bar{q}_{1} & D_{1} = \bar{q}_{1}^{2} - m^{2}, \\
& & D_{2} = (\bar{q}_{1} - k)^{2} - m^{2}, \\
& & D_{2} = (\bar{q}_{1} - k)^{2} - m^{2}, \\
& & D_{3} = \bar{q}_{2}^{2}, \\
& & D_{4} = (\bar{q}_{1} + \bar{q}_{2})^{2} - m^{2}
\end{array}$$

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& D_{3} = \bar{q}_{2}^{2}, \\
& D_{4} = (\bar{q}_{1} + \bar{q}_{2})^{2} - m^{2}
\end{array}$$

• Basis  $\{e_i\} \equiv \{k, k_{\perp}, e_3, e_4\}$  and coordinates  $\mathbf{z} = (x_1, x_2, x_3, x_4, y_1, y_2, y_3, y_4, \mu_{11}, \mu_{12}, \mu_{22})$  $\bar{q}_1^{(4\text{-dim})} = \sum_i x_i e_i, \qquad \bar{q}_2^{(4\text{-dim})} = \sum_i y_i e_i, \qquad (\bar{q}_i \cdot \bar{q}_j) = (\bar{q}_i^{(4\text{-dim})} \cdot \bar{q}_j^{(4\text{-dim})}) - \mu_{ij}$ 

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quotients

remainder

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\bar{q}_{2} & & D_{2} = (\bar{q}_{1} - k)^{2} - m^{2} , \\
\bar{q}_{2} & & D_{3} = \bar{q}_{2}^{2} , \\
D_{4} = (\bar{q}_{1} + \bar{q}_{2})^{2} - m^{2} \\
\end{array}$$

• Basis  $\{e_i\} \equiv \{k, k_{\perp}, e_3, e_4\}$  and coordinates  $\mathbf{z} = (x_1, x_2, x_3, x_4, y_1, y_2, y_3, y_4, \mu_{11}, \mu_{12}, \mu_{22})$  $\bar{q}_1^{(4-\text{dim})} = \sum_i x_i e_i, \qquad \bar{q}_2^{(4-\text{dim})} = \sum_i y_i e_i, \qquad (\bar{q}_i \cdot \bar{q}_j) = (\bar{q}_i^{(4-\text{dim})} \cdot \bar{q}_j^{(4-\text{dim})}) - \mu_{ij}$ • division of  $\mathcal{N}_{11234}$  modulo  $\mathcal{G}_{\mathcal{J}_{11234}} (= \mathcal{G}_{\mathcal{J}_{1234}})$ 

$$\mathcal{N}_{11234} = \underbrace{\mathcal{N}_{1234}D_1 + \mathcal{N}_{1134}D_2 + \mathcal{N}_{1124}D_3 + \mathcal{N}_{1123}D_4}_{\text{quotients}} + \underbrace{\Delta_{11234}}_{\text{remainder}}$$

• division of 
$$\mathcal{N}_{i_1i_2i_3i_4}$$
 modulo  $\mathcal{G}_{\mathcal{J}_{i_1i_2i_3i_4}}$  , e.g.

$$\mathcal{N}_{1234}/\mathcal{G}_{\mathcal{J}_{1234}} \Rightarrow \mathcal{N}_{1234} = \underbrace{\mathcal{Q}_{234}^{(1234)}D_1 + \mathcal{Q}_{134}^{(1234)}D_2 + \mathcal{Q}_{124}^{(1234)}D_3 + \mathcal{Q}_{123}^{(1234)}D_4}_{\text{quotients}} + \underbrace{\Delta_{1234}}_{\text{remainder}}$$

$$\mathcal{N}_{1134}/\mathcal{G}_{\mathcal{J}_{1134}} \Rightarrow \mathcal{N}_{1134} = \underbrace{\mathcal{Q}_{134}^{(1134)}D_1 + \mathcal{Q}_{114}^{(1134)}D_3 + \mathcal{Q}_{113}^{(1134)}D_4}_{\text{quotients}} + \underbrace{\Delta_{1134}}_{\text{remainder}}$$

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#### Divide-and-Conquer approach: a simple example

$$\begin{array}{cccc}
\bar{q}_{1} & D_{1} = \bar{q}_{1}^{2} - m^{2}, \\
& & D_{2} = (\bar{q}_{1} - k)^{2} - m^{2}, \\
& & D_{2} = (\bar{q}_{1} - k)^{2} - m^{2}, \\
& & D_{3} = \bar{q}_{2}^{2}, \\
& & D_{4} = (\bar{q}_{1} + \bar{q}_{2})^{2} - m^{2}
\end{array}$$

• Basis  $\{e_i\} \equiv \{k, k_{\perp}, e_3, e_4\}$  and coordinates  $\mathbf{z} = (x_1, x_2, x_3, x_4, y_1, y_2, y_3, y_4, \mu_{11}, \mu_{12}, \mu_{22})$  $\bar{q}_1^{(4\text{-dim})} = \sum_i x_i e_i, \qquad \bar{q}_2^{(4\text{-dim})} = \sum_i y_i e_i, \qquad (\bar{q}_i \cdot \bar{q}_j) = (\bar{q}_i^{(4\text{-dim})} \cdot \bar{q}_j^{(4\text{-dim})}) - \mu_{ij}$ • division of  $\mathcal{N}_{11234}$  modulo  $\mathcal{G}_{\mathcal{J}_{11234}} (= \mathcal{G}_{\mathcal{J}_{1234}})$  $\mathcal{N}_{11234} = \underbrace{\mathcal{N}_{1234}D_1 + \mathcal{N}_{1134}D_2 + \mathcal{N}_{1124}D_3 + \mathcal{N}_{1123}D_4}_{\mathcal{H}_1} + \underbrace{\Delta_{11234}}_{\mathcal{H}_1}$ 

quotients

remainder

22

• division of  $\mathcal{N}_{i_1i_2i_3i_4}$  modulo  $\mathcal{G}_{\mathcal{J}_{i_1i_2i_3i_4}}$ 

 $\mathcal{N}_{11234} = \underbrace{\mathcal{N}_{234}D_1^2 + \mathcal{N}_{134}D_1D_2 + \mathcal{N}_{124}D_1D_3 + \mathcal{N}_{123}D_1D_4 + \mathcal{N}_{114}D_2D_3 + \mathcal{N}_{113}D_2D_4}_{(sums of) quotients} + \underbrace{\Delta_{1234}D_1 + \Delta_{1134}D_2 + \Delta_{1124}D_3 + \Delta_{1123}D_4}_{remainders} + \Delta_{11234}$ 

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#### Divide-and-Conquer approach: a simple example

• after a further step (division  $N_{i_1i_2i_3}/\mathcal{G}_{\mathcal{J}_{i_1i_2i_3}}$ ) no quotient remains

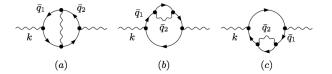
 $\mathcal{N}_{11234} = \Delta_{11234} + \Delta_{1234}D_1 + \Delta_{1134}D_2 + \Delta_{1124}D_3 + \Delta_{1123}D_4 + \Delta_{234}D_1^2 + \Delta_{114}D_2D_3 + \Delta_{113}D_2D_4$ 

the integrand decomposition becomes

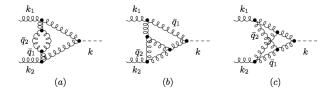
$$\begin{split} \mathcal{I}_{11234} &= \frac{\mathcal{N}_{11234}}{D_1^2 D_2 D_3 D_4} = \frac{\Delta_{11234}}{D_1^2 D_2 D_3 D_4} + \frac{\Delta_{1234}}{D_1 D_2 D_3 D_4} + \frac{\Delta_{1134}}{D_1^2 D_2 D_3} + \frac{\Delta_{1124}}{D_1^2 D_2 D_4} \\ &\quad + \frac{\Delta_{1123}}{D_1^2 D_2 D_3} + \frac{\Delta_{234}}{D_2 D_3 D_4} + \frac{\Delta_{114}}{D_1^2 D_4} + \frac{\Delta_{113}}{D_1^2 D_3} \\ \Delta_{11234} &= 16m^2 \left(k^2 + 2\,m^2 - k^2\epsilon\right) \,, \\ \Delta_{1234} &= 16 \left[ (q_2 \cdot k)(1 - \epsilon)^2 + m^2 \right] \,, \\ \Delta_{1124} &= -\Delta_{1123} = 8 \left(1 - \epsilon\right) \left[ k^2(1 - \epsilon) + 2\,m^2 \right] \,, \\ \Delta_{1134} &= -16m^2 \left(1 - \epsilon\right) \,, \\ \Delta_{113} &= -\Delta_{114} = \Delta_{234} = 8 \left(1 - \epsilon\right)^2 \,. \end{split}$$

#### Examples of divide-and-conquer approach

• Photon self-energy in massive QED,  $(4 - 2\epsilon)$ -dimensions



• Diagrams entering  $gg \rightarrow H$ , in  $(4 - 2\epsilon)$ -dimensions



#### **Conclusions and Outlook**

- Conclusions
  - We developed a general framework for the reduction at the integrand level
    - can be applied to any amplitude in any QFT, at every loop order
  - At one loop
    - naturally reproduces known results (OPP)
    - allows to express any amplitude in terms of known MIs
    - can be improved with the Laurent-expansion approach (NINJA)
  - At higher loops
    - it gives a recursive formula for the integrand decomposition
    - generates the form of the residue for every cut
    - can decompose any integrand with purely algebraic operations (divide-and-conquer approach)
- Outlook
  - application to a full 2-loop QED/QCD process
  - combine integrand reduction with other techniques (e.g. IBP)

# THANK YOU FOR YOUR ATTENTION

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